

# Mathematics and numerics for data assimilation and state estimation – Lecture 3



Summer semester 2020

## Summary of lecture 2

- Random vectors  $(X, Y) : \Omega \rightarrow A \times B$  and joint distributions

$$\mathbb{P}_{(X,Y)}((a, b)) = \mathbb{P}(X = a, Y = b).$$

- Independence of rv

$$\mathbb{P}(X = a, Y = b) = \mathbb{P}(X = a)\mathbb{P}(Y = b), \quad \forall a \in A \quad b \in B$$

and of events

$$\mathbb{P}(H_1 \cap H_2) = \mathbb{P}(H_1)\mathbb{P}(H_2)$$

- Expectation of  $X : \Omega \rightarrow A$ ,

$$\mu = \mathbb{E}[X] := \int_{\Omega} X(\omega)\mathbb{P}(d\omega) = \sum_{a \in A} a\mathbb{P}(X = a)$$

## Summary of lecture 2

- Variance of  $X$ . Defined for a scalar-valued rv (meaning  $X : \Omega \rightarrow A \subset \mathbb{R}^d$  with  $d = 1$ ),

$$\text{Var}(X) := \mathbb{E} [(X - \mu)^2].$$

- Property:  $\mu$  is the best constant-value approximation of  $X$  in the following sense

$$\mathbb{E} [(X - \mu)^2] \leq \mathbb{E} [(X - k)^2] \quad \text{for all } k \in \mathbb{R}.$$

## Plan for this lecture

- Conditional probabilities and expectations
- Conditioning on events: “probability of  $X$  given  $H$ ”

$$\mathbb{P}(X = a \mid H) \quad H \in \mathcal{F},$$

- Conditioning on rv: “probability of  $X$  given rv  $Y$ ”:

$$\mathbb{P}(X = a \mid Y)$$

Interesting property

$$\mathbb{E} [ |X - \mathbb{E}[X \mid Y]|^2 ] = \mathbb{E} [ |X - f(Y)|^2 ]$$

for any mapping  $f(Y) \in \mathbb{R}^d$ .

# Conditional probability

## Definition 1

For two events  $G, H \in \mathcal{F}$  where  $\mathbb{P}(H) > 0$ , the conditional probability of  $G$  given  $H$  is given by

$$\mathbb{P}(G | H) = \frac{\mathbb{P}(G \cap H)}{\mathbb{P}(H)}$$

Whenever  $\mathbb{P}(H) > 0$ , the mapping  $\mathbb{P}(\cdot | H) : \mathcal{F} \rightarrow [0, 1]$  is a probability measure.<sup>1</sup>

Verification:

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<sup>1</sup>And it remains to define  $\mathbb{P}(\cdot | H)$  for zero-probability events  $H$ .

Simplification in some settings (direct use of conditioning):

For  $X, Y$  and  $f(X, Y)$  discrete rv,

$$\mathbb{P}(f(X, Y) = c \mid Y = b) = \frac{\mathbb{P}(f(X, b) = c)}{\mathbb{P}(Y = b)}, \quad \text{if } \mathbb{P}(Y = b) > 0. \quad (1)$$

### Example 2

Let  $X_1, X_2, X_3 \sim \text{Bernoulli}(p)$  and independent rv. Let  $Z = X_1 + X_2 + X_3$ .  
Compute

$$\mathbb{P}(Z \geq 1 \mid X_1 = 0)$$

**Solution:**

### Example 3 (Example where conditioning information is used "implicitly")

Let  $X_1, X_2, X_3 \sim \text{Bernoulli}(p)$  and independent rv. Let  $Z = X_1 + X_2 + X_3$ .  
Compute

$$\mathbb{P}(X_1 = 1 \mid Z = 2)$$

**Solution:**

#### Definition 4 (Conditional expectation)

For a discrete rv  $X : \Omega \rightarrow A$  and an event  $H \in \mathcal{F}$  with  $\mathbb{P}(H) > 0$ , we define the conditional expectation of  $X$  given  $H$  as

$$\mathbb{E}[X | H] := \int_{\Omega} X(\omega) \mathbb{P}(d\omega | H) = \sum_{a \in A} a \mathbb{P}(X = a | H)$$

■ Property:

$$\mathbb{E}[X | H] = \mathbb{E}[X \mathbf{1}_H] / \mathbb{P}(H) \quad (2)$$

Verification:

■ Implication:  $\mathbb{E}[|X| | H] \leq \mathbb{E}[|X|] / \mathbb{P}(H)$ .



## Example 5

Let  $X$  be a three-sided fair die, meaning

$$\mathbb{P}(X = k) = \frac{1}{3} \quad \text{for } k = 1, 2, 3.$$

Compute  $\mathbb{E}[X \mid X \geq 2]$ .

**Solution:**

## Conditioning on zero-probability events

For events  $G, H \in \mathcal{F}$ , it is not clear how interpret the definition

$$\mathbb{P}(G | H) := \frac{\mathbb{P}(G \cap H)}{\mathbb{P}(H)}$$

when  $\mathbb{P}(H) = 0$ .

**Is an extension of the definition needed?** May not seem needed as zero-probability events “never” happen anyway, but often it is convenient to use the same co-domain for all rv studied, say for example

$$X_k : \Omega \rightarrow \mathbb{N}$$

with  $X_k(\omega) = k$  for  $\omega \in \Omega$  and  $X_k(\omega) = 0$  for  $\omega \in \Omega \setminus \{k\}$  for  $k = 1, 2, \dots$

**Also** any event  $\{Y = y\}$  is a zero-probability event for a continuous rv!

## Conditioning on zero-probability events 2

### Definition 6 (Division-by-zero convention)

For any  $c \in \mathbb{R}$  we will, in all of this course, make use of the following convention

$$\frac{c}{0} := 0.$$

**Motivation:** Then  $\frac{a}{b}$  is defined for any  $a, b \in \mathbb{R}$ , but it gives algebra a quirk

$$b(a/b) = \begin{cases} a & \text{if } b \neq 0 \text{ or } a = 0 \\ 0 & \text{if } b = 0. \end{cases}$$

## Definition 7 (Generalization of Definition 1)

For **any** pair of events  $G, H \in \mathcal{F}$ , we define

$$\mathbb{P}(G | H) := \frac{\mathbb{P}(G \cap H)}{\mathbb{P}(H)}$$

where we note that by the division-by-zero convention

$$\mathbb{P}(G | H) = 0 \quad \text{if } \mathbb{P}(H) = 0.$$

### Implications:

- The definition of conditional expectation “naturally” extends to any zero-probability events  $H \in \mathcal{F}$ :

$$\mathbb{E}[X|H] := \sum_{a \in A} a \mathbb{P}(X = a | H) = 0.$$

- Direct use of conditioning, cf. equation (1), extends. Meaning,

$$\mathbb{P}(f(X, Y) = c | Y = b) = \frac{\mathbb{P}(f(X, b) = c)}{\mathbb{P}(Y = b)}, \quad \text{also if } \mathbb{P}(Y = b) = 0.$$

## Conditioning on random variables

- We have defined the conditional probability  $\mathbb{P}(G | H)$  for any pair events  $G, H$ .
- So for rv  $X : \Omega \rightarrow A$  and  $Y : \Omega \rightarrow B$ , the following quantities are all defined

$$\mathbb{P}(X = a | Y = b) \quad \text{for any } a \in A, b \in B.$$

- Fixing the event  $\{X = a\}$ , we may introduce the function  $\psi : B \rightarrow [0, 1]$

$$\psi(b) = \mathbb{P}(G | \{Y = b\})$$

- and the function  $\phi : \Omega \rightarrow [0, 1]$  by

$$\phi(\omega) := \mathbb{P}(X = a | \{Y = Y(\omega)\})$$

( curly brackets in the  $\{Y = Y(\omega)\}$  notation here is only used to emphasize that we have events and is not really needed).

## Conditioning on random variables 2

- The mapping  $\phi(\omega)$  was introduced to clarify that  $\mathbb{P}(X = a \mid \{Y = Y(\omega)\})$  is a function of  $\omega$ .
- The customary notation for these conditional probabilities is as follows:

### Definition 8 (Probability of $X$ given $Y$ )

Consider the discrete rv  $X$  and  $Y$  on the previous slide. Then for each  $a \in A$ , the mapping  $\mathbb{P}(X = a \mid Y) : \Omega \rightarrow [0, 1]$  is the discrete rv defined by

$$\mathbb{P}(X = a \mid Y)(\omega) = \mathbb{P}(X = a \mid \{Y = Y(\omega)\}).$$

Verification that  $\phi(\omega) = \mathbb{P}(X = a | Y)(\omega)$  is a discrete rv:

- The set of outcomes/ image space

$$\begin{aligned}\phi(\Omega) &= \cup_{\omega \in \Omega} \{\mathbb{P}(X = a | Y = Y(\omega))\} \\ &= \cup_{b \in B} \{\mathbb{P}(X = a | Y = b)\} =: C \subset [0, 1]\end{aligned}$$

is countable since  $B$  is countable.

- For each  $c \in C$ , there exists a  $b(c) \in B$  such that

$$c = \mathbb{P}(X = a | Y = b(c))$$

and

$$\phi^{-1}(c) = \{\omega \in \Omega | Y(\omega) = b(c)\} \in \mathcal{F}.$$

## Example 9

Consider the a coin toss  $X : \Omega \rightarrow \{0, 1\}$  and a die roll  $Y : \Omega \rightarrow \{1, 2, 3\}$ , on sample space  $\Omega = \{Heads, Nose, Tails\}$  with

$$X^{-1}(1) = \{Heads, Nose\} \quad \text{and} \quad X^{-1}(0) = \{Tails\}$$

$$Y^{-1}(1) = \{Heads\}, \quad Y^{-1}(2) = \{Nose\} \quad \text{and} \quad Y^{-1}(3) = \{Tails\}.$$

and

$$\mathbb{P}(Heads) = \mathbb{P}(Nose) = 1/4, \quad \text{and} \quad \mathbb{P}(Tails) = 1/2.$$

Then

$$\mathbb{P}(X = 0 \mid Y)(Heads) =$$

$$\mathbb{P}(Y = 1 \mid X)(Nose) =$$



### Definition 10 (Expectation of $X$ given $Y$ )

For discrete rv  $X : \Omega \rightarrow A \subset \mathbb{R}^d$  and  $Y : \Omega \rightarrow B \subset \mathbb{R}^k$  with  $|\mathbb{E}[X]| < \infty$ , the mapping  $\mathbb{E}[X | Y] : \Omega \rightarrow \mathbb{R}^d$  is defined

$$\mathbb{E}[X | Y](\omega) := \sum_{a \in A} a \mathbb{P}(X = a | Y)(\omega) = \sum_{a \in A} a \mathbb{P}(X = a | \{Y = Y(\omega)\}).$$

Note,  $\mathbb{E}[X | Y]$  is a discrete rv.

### Example 11

Consider the Bernoulli rv  $X, Y$  with joint probabilities

$$\mathbb{P}(X = i, Y = j) = \begin{bmatrix} 1/8 & 1/4 \\ 1/2 & 1/8 \end{bmatrix} \quad i, j \in \{0, 1\}$$

$$\mathbb{E}[Y | X](\text{Heads}) =$$

### Definition 10 (Expectation of $X$ given $Y$ )

For discrete rv  $X : \Omega \rightarrow A \subset \mathbb{R}^d$  and  $Y : \Omega \rightarrow B \subset \mathbb{R}^k$  with  $|\mathbb{E}[X]| < \infty$ , the mapping  $\mathbb{E}[X | Y] : \Omega \rightarrow \mathbb{R}^d$  is defined

$$\mathbb{E}[X | Y](\omega) := \sum_{a \in A} a \mathbb{P}(X = a | Y)(\omega) = \sum_{a \in A} a \mathbb{P}(X = a | \{Y = Y(\omega)\}).$$

Note,  $\mathbb{E}[X | Y]$  is a discrete rv.

Note that,

$$\mathbb{E}[X | Y](\omega) = \mathbb{E}[X | \{Y = Y(\omega)\}]$$

and that it can be associated to a deterministic mapping  $g : \mathbb{R}^k \rightarrow \mathbb{R}^d$  as follows

$$g(Y(\omega)) = \mathbb{E}[X | Y = Y(\omega)]. \quad (3)$$

## Motivation for $\mathbb{E}[X | Y]$

Say you have an observation  $Y(\omega)$  (i.e., you know  $Y(\omega)$  but not  $\omega$ ), and that what you really seek is the value of  $X(\omega)$ . Then what is the best function  $g(Y(\omega))$  to approximate  $X(\omega)$ ?

### Theorem 11 (Mean-square sense best approximation)

For discrete rv  $X : \Omega \rightarrow A \subset \mathbb{R}^d$  and  $Y : \Omega \rightarrow B \subset \mathbb{R}^k$  with  $\mathbb{E}[X^2] < \infty$ , it holds that

$$\mathbb{E}[|X - \mathbb{E}[X | Y]|^2] \leq \mathbb{E}[|X - f(Y)|^2]$$

for all  $f : \mathbb{R}^k \rightarrow \mathbb{R}^d$  such that  $\mathbb{E}[|f(Y)|^2] < \infty$ .

**Interpretation** Since the constant function  $f(Y) = \mathbb{E}[X]$  is one possible mapping, we conclude that

$$\mathbb{E}[|X - \mathbb{E}[X | Y]|^2] \leq \mathbb{E}[|X - \mathbb{E}[X]|^2].$$

To prove Theorem 11, we will need a few intermediary results.

### Lemma 12 (The tower property )

For discrete rv  $X : \Omega \rightarrow A \subset \mathbb{R}^d$  and  $Y : \Omega \rightarrow B \subset \mathbb{R}^k$  with  $|\mathbb{E}[X]| < \infty$ , it holds that

$$\mathbb{E}[\mathbb{E}[X | Y]] = \mathbb{E}[X].$$

**Proof:**

### Lemma 13 (The **Direct conditioning of expectations** )

For the setting in Lemma 12, it holds for any mapping  $f : \mathbb{R}^d \times \mathbb{R}^k \rightarrow \mathbb{R}$  such that  $|\mathbb{E}[f(X, Y)]| < \infty$  that

$$\mathbb{E}[f(X, Y) \mid Y = b] = \mathbb{E}[f(X, b) \mid Y = b] \quad \forall b \in B.$$

**Special case:**  $f(x, y) = g(x)h(y)$  yields

$$\mathbb{E}[g(X)h(Y) \mid Y = b] = h(b)\mathbb{E}[g(X) \mid Y = b] \quad \forall b \in B$$

Since this holds for all  $b$ ,

$$\mathbb{E}[g(X)h(Y) \mid Y = b] = h(Y)\mathbb{E}[g(X) \mid Y].$$

**And tower property 2:**

$$\mathbb{E}[h(Y)\mathbb{E}[g(X) \mid Y]] = \mathbb{E}[h(Y)g(X)] \quad (4)$$

Using (4), let us prove Theorem 11 in the 1D setting, i.e., that

$$\mathbb{E}[(X - \mathbb{E}[X | Y])^2] \leq \mathbb{E}[(X - f(Y))^2]$$

for all  $f : \mathbb{R} \rightarrow \mathbb{R}$  with  $\mathbb{E}[(f(Y))^2] < \infty$ .

**Proof:**

$$\begin{aligned} \mathbb{E}[(X - f(Y))^2] &= \mathbb{E}\left[\left((X - \mathbb{E}[X | Y]) + (\mathbb{E}[X | Y] - f(Y))\right)^2\right] \\ &= \end{aligned}$$

For  $X : \Omega \rightarrow A \subset \mathbb{R}^d$  and  $Y : \Omega \rightarrow B$ , the mapping

$$g(b) := \mathbb{E}[X \mid Y = b]$$

satisfies

$$g(Y(\omega)) := \mathbb{E}[X \mid Y = Y(\omega)].$$

**Conclusion:**  $\mathbb{E}[X \mid Y]$  is an rv induced from the rv  $Y$  through the mapping  $g$ .

**Question:** Is  $\mathbb{E}[X \mid Y]$  in some sense unique?

**Question:** Given a candidate mapping  $g : B \rightarrow \mathbb{R}^d$ , is there a way to verify whether  $g(Y) = \mathbb{E}[Y \mid X]$ ?

### Definition 14 ( $\mathbb{P}$ -almost surely equal)

Two rv  $X, Y$  are said to be  $\mathbb{P}$ -almost surely equal provided

$$\mathbb{P}(\{\omega \in \Omega \mid X(\omega) = Y(\omega)\}) = 1.$$

We write

$$X = Y \quad \mathbb{P} - a.s.$$

(or just “a.s.” whenever it is clear which probability measure  $\mathbb{P}$  is considered).

Motivation:

### Example 15

$X : \Omega \rightarrow \{0, 1\}$  and  $Y : \Omega \rightarrow \{0, 1, 2\}$  with

$$\mathbb{P}(X = Y) = 1 \quad \text{and} \quad \{Y = 2\} \neq \emptyset.$$

Then  $X(\omega) \neq Y(\omega)$  for any  $\omega \in \{Y = 2\}$ , but  $X = Y$  a.s.



## Theorem 16

Consider the setting in Lemma 12.

If  $g : \mathbb{R}^k \rightarrow \mathbb{R}^d$  is a mapping such that for every bounded mapping  $f : \mathbb{R}^k \rightarrow \mathbb{R}$ ,

$$\mathbb{E}[f(Y)g(Y)] = \mathbb{E}[f(Y)X] \quad (5)$$

then

$$g(Y) = \mathbb{E}[X | Y] \quad \text{a.s.}$$

**Interpretation:**  $\mathbb{E}[X | Y]$  is a a.s. unique rv of form  $g(Y)$  satisfying (5).

**Usage:** If a mapping  $B \ni b \mapsto g(b) \in \mathbb{R}^d$  satisfies (5), i.e.,

$$\sum_{b \in B} f(b)g(b)P(Y = b) = \sum_{a \in A, b \in B} f(b)aP(X = a, Y = b) \quad \forall f : B \rightarrow \mathbb{R},$$

then  $g(Y(\omega)) = \mathbb{E}[X|Y](\omega)$  for  $\mathbb{P}$ -almost all  $\omega \in \Omega$ .

## Next time

- Convergence of random variables
- Random walks and discrete time Markov Chains